

## Day 2: Problem Analysis

15.04.2015

# Problem A. Another 2048 Problem

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# Problem statement

- You are given multiset of integers
- In one turn you can take two equal numbers, remove these two numbers from the set and insert their sum instead
- You have to find the number of subsets, such that you can somehow get 2048 making several turns

# Solution

- Sum up only powers of two, otherwise can't get 2048 only multiplying by 2
- Remove all  $m$  non-powers of two, and multiply the answer by  $2^m$
- Numbers don't exceed 2047, all the subsets with sum of their numbers  $\geq 2048$  are good
- Calculate number of bad subsets  $\Rightarrow$  knapsack problem
- $f[i][s]$  — the number of subsets, containing numbers up to  $2^i$  with sum equal to  $s$
- Optimize: one only needs  $f[i][s]$ , such that  $s$  is multiple of  $2^i$ , otherwise round down

## Problem B. Big Kingdom

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# Problem statement

- Given  $n$  guards in points  $(x_i, y_i)$  on the plane, and velocity for every guard  $v_i$
- The point is guarded by the  $i$ -th guard, if the time to get to this points is strictly less, than for any other guard
- Find if the area guarded by every guard is infinite

# Solution

- If for some guard there exists faster than him, then his area is not infinite
- Get all the guards, build convex hull
- Guard on the border  $\Leftrightarrow$  his guarded area is infinite

## Problem C. Construct The Array

# Problem C. Construct The Array



# Problem statement

- You have an array. Two kinds of operations made:
  - 1 Add  $v$  to all  $a_x$ , such that  $\gcd(x, n) = d$
  - 2 Calculate  $\sum_{i=1}^x a_i$

## Solution

- Suppose  $d = 1$ ,  $n = p_1^{w_1} \cdot p_2^{w_2} \cdot \dots \cdot p_k^{w_k}$ , then  $x$  shouldn't be divisible to all  $p_i$
- Use inclusion-exclusion principle, add  $(-1)^{|A|} \cdot v$  to  $f \left[ \prod_{i \in A} p_i \right]$
- $\sum_{i=1}^x a_i = \sum_{i=1}^x f[i] \cdot \lfloor \frac{x}{i} \rfloor$
- $d \neq 1$  is almost the same, supposing  $n := \frac{n}{d}$  and  $x = d \cdot \prod_{i \in A} p_i$
- Calculate  $\sum_{i=1}^x f[i] \cdot \lfloor \frac{x}{i} \rfloor$  in  $O(\sqrt{x})$ , since there are about  $\sqrt{x}$  different values of  $\lfloor \frac{x}{i} \rfloor$

# Problem D. Dictator

## Problem D. Dictator

# Problem statement

- Given a tournament
- You need to arrange the vertices in order, such that for every previous vertex there is a path of length 1 or 2, passing only previous vertices

# Solution

- There is always the answer
- Look at the vertex  $v$ , that has the most incoming edges  $k$ , make it last in the order
- Need to prove, that from every other vertex  $u$  there is the path of length 1 or 2
- Two cases:
  - $uv \in E$ , then the edge is the path
  - Otherwise, there exists  $w$ , such that  $uw \in E$  and  $wv \in E$ .  
Suppose it's not, then  $u$  has  $k$  incoming edges from first type of vertices and one incoming edge from  $v$ , making in total  $k + 1$  edges. So  $k$  is not maximum number of incoming edges, leading to the contradiction

## Problem E. Electricity and Magic

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# Problem statement

- Given a grid, consisting of 0s and 1s
- You can make two types of operations:
  - 1 Flip bits that are adjacent to  $(i, j)$
  - 2 Flip bits that are adjacent to  $(i, j)$  and  $(i, j)$  itself
- Find minimum number of moves needed to obtain a field filled with 0s

# Solution

- Dynamic programming approach
- For every cell you need to remember one of three states:
  - You made moves on cell  $(i, j)$
  - You didn't make any move on cell, it is equal to 0
  - You didn't make any move on cell, it is equal to 1
- If you made move on cell, then you can make it either 0 or 1, changing the type of one of the moves
- $O(nm3^m)$  solution



# Problem F. Fight

## Problem F. Fight

# Problem statement

- Hero attacks monster reducing monster's HP  $h$  by  $a$
- Monster HP increases by  $b$  every second, after the attack if it's been made
- Hero can't make  $k + 1$  consecutive attacks
- Can hero kill the monster?

# Solution

- Killing in first round  $\Leftrightarrow h \leq a$
- Killing before rest  $\Leftrightarrow h + b(k - 1) \leq ak$
- Killing with rest  $\Rightarrow b(k + 1) < ak$

## Problem G. Go and restore!

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# Problem statement

- Given multiplication table base  $p$
- Digits were shuffled, and table was rearranged
- Find the permutation, that shuffles the digits

# Solution

- Look at row which corresponds to  $p - 1$
- It looks like  $[0][0]$ ,  $[0][p - 1]$ ,  $[1][p - 2]$ ,  $[2][p - 3] \dots [p - 2][1]$
- No other row has  $p - 2$  as highest digit, so it is the only row, that has  $p - 1$  different highest digits
- Find that row, learn what  $p - 1$  maps to
- Using the fact that  $(p - 1) \times k = [k - 1][p - k]$ , find  $map[k - 1]$  from  $map[k]$

# Problem H. Hidden Integer

## Problem H. Hidden Integer

# Problem statement

- Given integer  $x$
- You make  $k$  moves
- During  $i$ -th move, you replace  $x$  by minimal  $y$ , such that  $y \geq x$  and  $y$  is divisible by  $i$
- Find the value of  $x$  after  $k$  moves



# Solution

- Let's represent the value of  $x$  after  $i$  moves as  $a \cdot i$
- After the move  $x$  will equal to  $b \cdot (i + 1)$  for some  $b$ , such that  $b \cdot (i + 1) \geq a \cdot i$
- So  $b = \left\lceil \frac{a \cdot i}{i + 1} \right\rceil$
- It can be proven, that if  $a$  stopped changing, then it never changes again, and this moment happens fast

# Problem J. Just do it

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# Problem statement

- You are to calculate permanent of matrix... Of  $m$  matrices.
- $i$ -th of them:
  - If  $x \neq y$ , then  $w_{x,y} = 1$
  - If  $x = y$  and  $x \leq n$ , then  $w_{x,y} = a_x$
  - If  $x = y$  and  $x > n$ , then  $w_{x,y} = 0$

# Solution

- Let's solve for first matrix
- The permanent of this matrix equals to:

$$\sum_{k=0}^n \left( \sum_{|A|=k} \prod_{x \in A} a_x \right) \cdot f_{n-k}$$

- Where  $f_x$  — is the number of permutations of length  $x$  without fixed points
- Generally, the permanent for  $i$ -th matrix is

$$\sum_{k=0}^{n+i-1} \left( \sum_{|A|=k} \prod_{x \in A} a_x \right) \cdot f_{n+i-1-k}$$

- It's easy to see:  $\sum_{k=0}^n \left( \sum_{|A|=k} \prod_{j \in A} a_j \right) x^k = \prod_{i=1}^n (1 + a_i x)$

# Solution

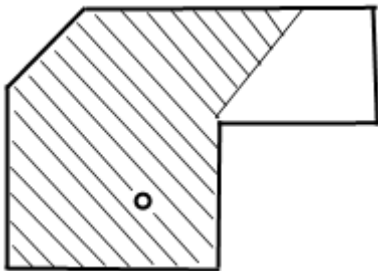
- To calculate  $\prod_{i=1}^n (1 + a_i x)$  use Fast Fourier Transform and Divide and Conquer technique
- Calculate  $f_n = \sum_{i=0}^n \frac{(-1)^i n!}{i!}$
- Computing permanents for all matrices is also reduced to polynomial multiplication
- Solution time complexity:  $O(n \log^2 n)$

# Problem L. Light Source

## Problem L. Light Source

# Problem statement

- You are given polygon and a point inside.
- Calculate the area visible from this point



# Solution

- Use scanline technique rotating the ray from the given point
- Maintain the set of sides of polygon, that intersect this ray, sorted by distance between center and intersection
- When vertex is met, either the segment appears or disappears
- For angle between two events, add the triangle area to the answer