Combinatorics compilation contest Problem analysis

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- You are given two sequences a_i and b_i and a prime modulo m.
- You have to calculate the number of pairs *i* and *j* such that $a_i \cdot b_j \mod m \le L$.

- Let us find a generator g for modulo m.
- Represent a_i and b_i as g^{α_i} and g^{β_i} , i.e., by calculating how many from the sequences are equal g^i modulo m.
- Multiplication becomes addition: $a_i \cdot b_j = g^{\alpha_i + \beta_j}$
- Consider a_i and b_j as polynomials, calculate the convolution using FFT.
- After convolution go back to powers of g and get the result.

Calculate the number of strings with hash code equal to x modulo $998244353 = 7 \cdot 17 \cdot 2^{23} + 1$.

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 dp[i][j] — the number of strings with length i that have hash j modulo m.

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$$dp[a+b][j] = \sum_{l=0}^{m} dp[a][l \cdot p^{-b}] \cdot dp[b][j-l].$$

- Consider dp[i](x) as the polynomial with coefficients dp[i][j].
- Equation above looks like the multiplication of dp[b](x) and slightly modified dp[a](x) (dp[a] → dp[a · p^b]).
- We could calculate all $dp[2^k](x)$ in $O(m \log m \log n)$ using FFT.
- Then we have to multiply $dp[2^k](x)$ corresponding to ones in *n*'s binary representation.

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- You are given 500 integers not bigger than 10^{12} .
- You have to find two distinct subsets with the same sum.

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- Leave only 50 integers (any of them).
- Since $2^{50} > 50 \cdot 10^{12}$ there exist two equal subsets.
- Take random 6 integers out of chosen 50.
- Add them up and store in hashset.
- Repeat until you find the answer.

- 25 times perform the following operation.
- Take $7 \cdot 10^6$ subsets of size 3.
- Sort them and take the two with the smallest difference (this difference is less than $\frac{3 \cdot 10^{12}}{7 \cdot 10^6} < 5 \cdot 10^5$).
- Remove these 6 values from the initial set.
- Now, we have 25 differences less than $5 \cdot 10^5$. Since, $25 \cdot 5 \cdot 10^5 < 2^{25}$, there exist subsets with equal sum.

- You are given the string.
- Run is a maximal sequence of contiguous equal characters.
- You have to calculate the number of strings with the same characters as given string and with the same number of runs.

- We will add symbols from 'a' to 'z' and calculate dp[c][r] the number of strings with the number of runs r after adding c first symbols.
- Suppose, we know the answer for first c 1 symbols and we add symbol c.
- Let us insert into string with r runs. There are two different types of insert positions: r + 1 "in-between", giving 1 additional run, and M - r "in-run", giving 2 additional runs.
- Fix a "in-between" and b "in-run". We add to $dp[c][r + a + 2 \cdot b]$ value $\binom{r+1}{a} \cdot \binom{M-r}{b} \cdot \binom{N_c-1}{a+b-1} \cdot dp[c-1][r]$.

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- There is a game. Each step: $(A, B) \rightarrow (A k \cdot B, B)$ or $(A, B) \rightarrow (A, B k \cdot A)$.
- The first one who gets zero or below loses.

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- Implement a naive solution, notice that all states (A, B)A ≥ B are winning if A ≥ B(A).
- B(A) is also increasing.
- When $A \ge 2 \cdot B$ (A, B) is a winning position, because we could get to $(A \mod B, B)$ or $(A \mod B + B, B)$.
- At least one of those positions is losing.
- Otherwise, you can check if it's winning or losing recursively in O(log(n)) steps.
- Use two pointers technique to determine B(A).
- The answer could be calculated simply by iterating though all A in $[A_1, A_2]$.

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- Number in S is considered **pure** if starting from it, you can continue taking its rank in S, and get a number in S, until in finite steps you hit 1.
- Find the number of ways to pick S with pure n from $\{1, \ldots, n\}$.

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• Let us calculate dp[n][k] — the number of subsets of size k where the biggest value is n and it is pure.

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$$dp[n][k] = \sum_{i=1}^{k-1} {n-k-1 \choose k-i-1} \cdot dp[k][i].$$

• Note, that you have to calculate dp one time for all the test cases.

- Cryptarithm equation is an addition equation in which all digits in one column are different.
- You need to calculate the number of cryptarithm equations with sum *N* in the given base *B*.

- First, calculate an auxiliary DP:
- count[k][v][f] the number of ways to represent v as a sum of k different non-negative numbers less than B. f is true, is there's a zero among those K numbers.
- Note that $v \leq B^2$.
- The main solution builds cryptarithm equations starting from the least significant digits.

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- dp[i][v][k][F] the number of cryptarithm equations on the last i digits, such that its sum conicides with i last digits of N, carry is v, is a from filled columns, k is a number of non-finished rows and F is a boolean flag indicating is there was a zero in last column.
- The set of these parameters is enough to make a DP transition.
- For the next column we could choose the number of rows *r* that will be continued, how much we will add to get *N_i*, and if this column contains 0 or not.
- Multiply by some binomial coefficients and factorials.

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- Fill a rectangular grid with letters.
- Letters in each row and column should increase.

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- Consider the area containing all symbols $\leq c$
- It is a **Young diagram** for every c
- There are $\binom{2n}{n} \sim \frac{4^n}{\sqrt{n}}$ monotone paths from (0, n) to (n, 0)



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- *dp*[*P*][*c*][*k*] is the number of ways to fill a Young diagram *P* with symbols up to *c* and *c* only appears in columns ≤ *k*.
- Add one more column and recalculate the path P.
- Generate all monotone paths before calculating DP and map them into integers.

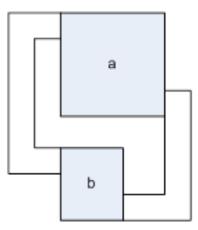
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You have to build a grid such that the number of down-right paths from upper-left corner to lower-right corner is k.

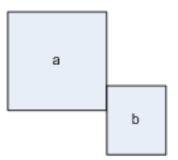
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I. The Reverse Problem About The Turtle Solution

The number of paths a and b to get a + b.



The number of paths a and b to get $a \cdot b$.



I. The Reverse Problem About The Turtle Solution

- We get the answers for powers of 2 applying the first constuction.
- Decompose N on a powers of 2 and apply the second construction.
- Example for 19.

